

The Phantom of Differential Characteristics

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Motivation

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DISTINGUISHER +

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For various application scenarios, we often assume the ability of an attacker to control the keys:

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For various application scenarios, we often assume the ability of an attacker to control the keys:

- Single-key model
- Open-key model
 - ▶ related-key attack
 - ▶ weak-key attack
 - ▶ known-key attack

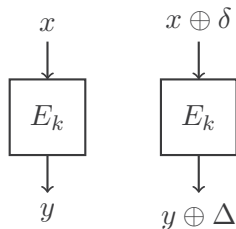
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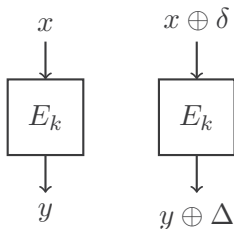
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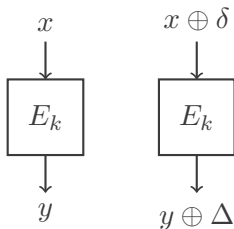
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- Track probabilistic difference propagation



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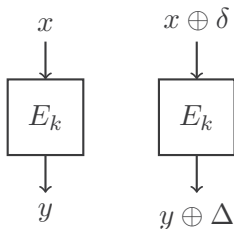
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Differential cryptanalysis

- One of the most extensively studied cryptanalytic techniques
- Track probabilistic difference propagation
- Differential characteristics and differentials
- Distinguish from random and key recovery



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Assumptions

- Markov cipher
- Independently random round keys
- Hypothesis of stochastic equivalence

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However, an attacker targets on one secret key.

- The probability of a differential distinguisher determines the attack complexity
- Differential or impossible differential?

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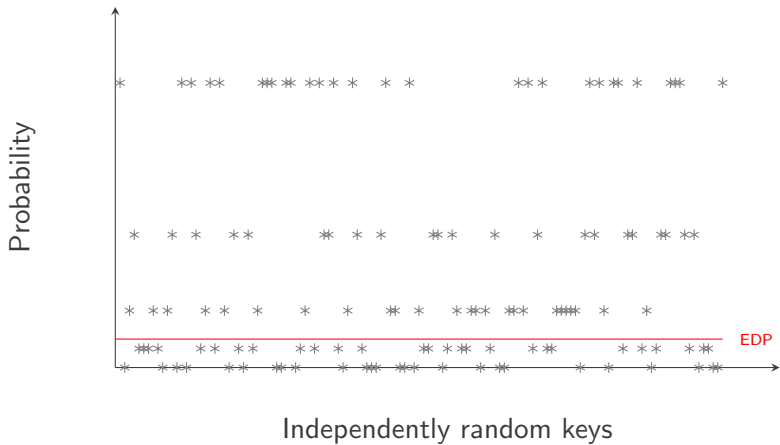
[Leu12] G. Leurent. Analysis of differential attacks in ARX constructions. ASIACRYPT 2012

[KNP+15] D. Khovratovich, I. Nikolić, J. Pieprzyk, P. Sokołowski, R. Steinfeld. Rotational cryptanalysis of ARX revisited. FSE 2015

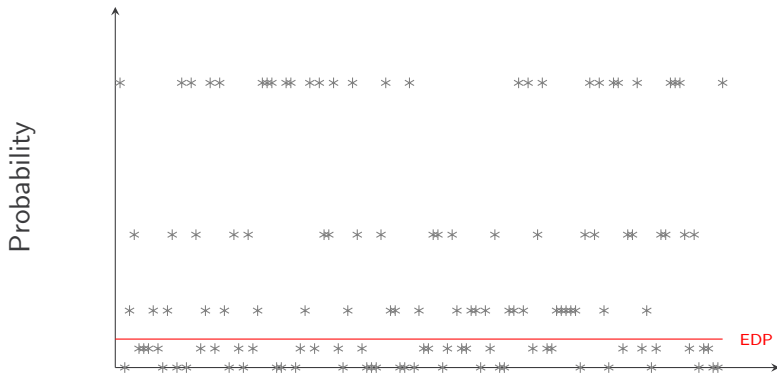
[DR07] J. Daemen, V. Rijmen. Plateau characteristics. IET information security, 2007

[CLN+17] A. Canteaut, E. Lambooj, S. Neves, S. Rasoolzadeh, Y. Sasaki, M. Stevens. Refined Probability of Differential Characteristics Including Dependency Between Multiple Rounds. IACR ToSC 2017 (2)

Motivation



Motivation



Independently random keys

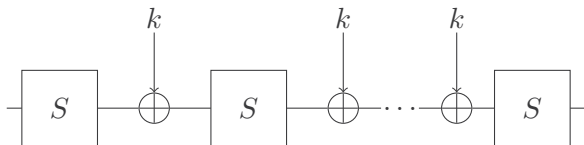
To what extent can we rely on the Assumptions?

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Enumerate characteristics under the Assumptions:

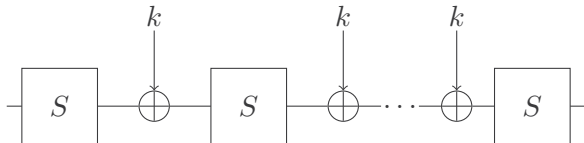
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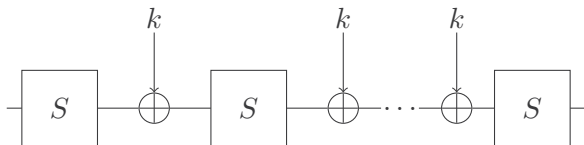
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- For a fixed key, # characteristics = 2^{15}

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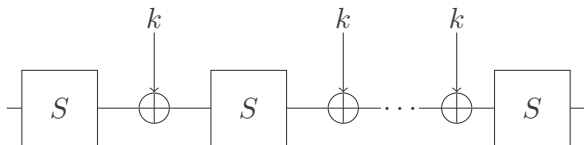
Enumerate characteristics under the Assumptions:



- For a fixed key, # characteristics = 2^{15}
- Under the Assumptions, # characteristics = $2^8 \times 2^7 \times \dots \times 2^7 = 2^{7r+8}$

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Enumerate characteristics under the Assumptions:



- For a fixed key, # characteristics = 2^{15}
- Under the Assumptions, # characteristics = $2^8 \times 2^7 \times \dots \times 2^7 = 2^{7r+8}$
- A characteristic generated under the Assumptions is “almost” impossible in reality.

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To study differential probability in fixed-key block ciphers and permutations

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It is crucial to ask:

- $EDP \neq 0$ while $DP = 0$ for all keys?
- Differential characteristics enumeration?
- Characteristics-based attacks?
- Compute DP under any given key?
- Design better key schedules and/or constants?

Effective Keys and Singular Characteristics

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A key is effective for a characteristic if the characteristic is of nonzero probability under the key.

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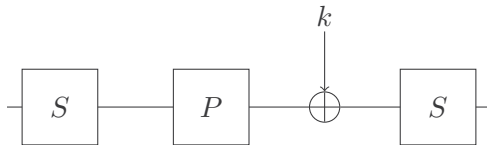
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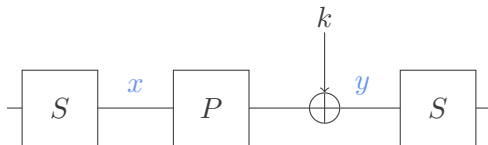
If no effective key exists, it is called a *singular characteristic*.

Effective Keys



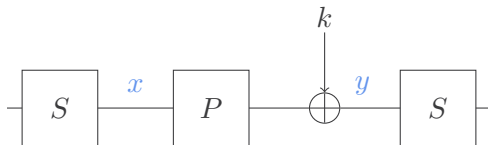
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- SPN cipher with keys XORed after the linear layer
- Right output and right input of the Sboxes
- Effective key candidates: $k = Px \oplus y$

Singular Characteristics

$$\alpha_0 \xrightarrow{S} \beta_0 \xrightarrow{P} \alpha_1 \xrightarrow{S} \beta_1 \xrightarrow{P} \alpha_2 \xrightarrow{S} \beta_2 \xrightarrow{P} \alpha_3 \xrightarrow{S} \beta_3 \xrightarrow{P} \alpha_4$$

Singular Characteristics

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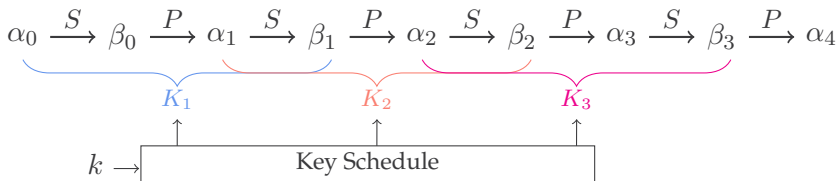
A blue curly brace is drawn under the first four terms of the sequence, $\alpha_0, \beta_0, \alpha_1, \beta_1$. Below the center of the brace is the label K_1 .

Singular Characteristics

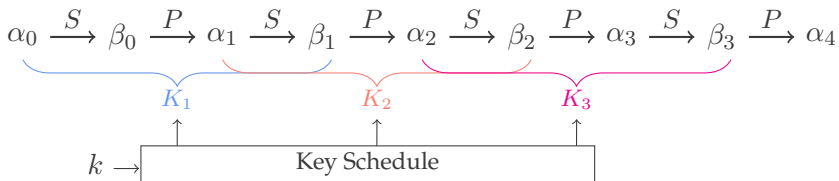
$$\alpha_0 \xrightarrow{S} \beta_0 \xrightarrow{P} \alpha_1 \xrightarrow{S} \beta_1 \xrightarrow{P} \alpha_2 \xrightarrow{S} \beta_2 \xrightarrow{P} \alpha_3 \xrightarrow{S} \beta_3 \xrightarrow{P} \alpha_4$$

K_1 K_2 K_3

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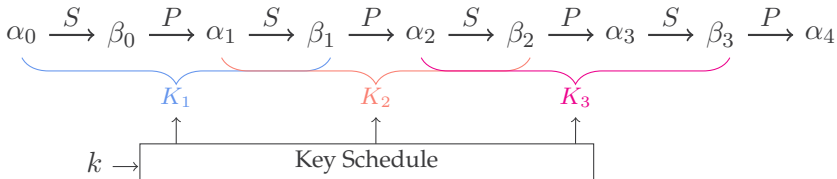


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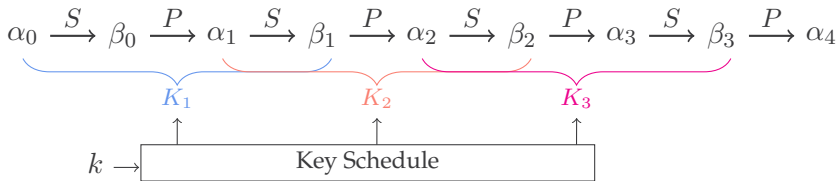
- When the difference propagation is legal, the effective key set of a 2-round characteristic is non-empty.

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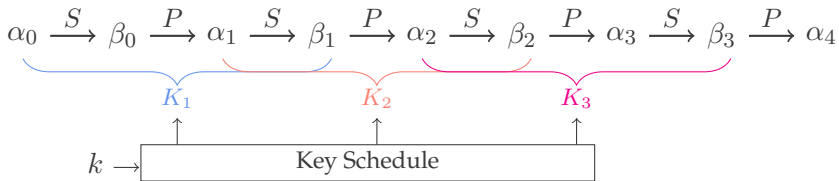
- When the difference propagation is legal, the effective key set of a 2-round characteristic is non-empty.
- Effective keys derived from two consecutive rounds may not be compatible with the key schedule.

Singular Characteristics



Procedure:

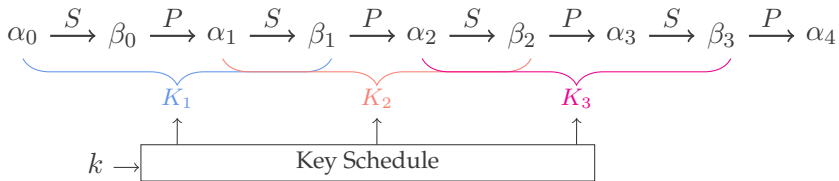
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Procedure:

1. Conditions on K_i to be effective

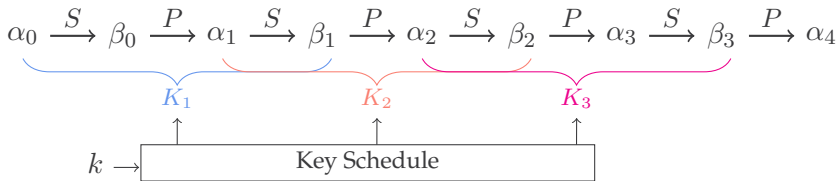
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Procedure:

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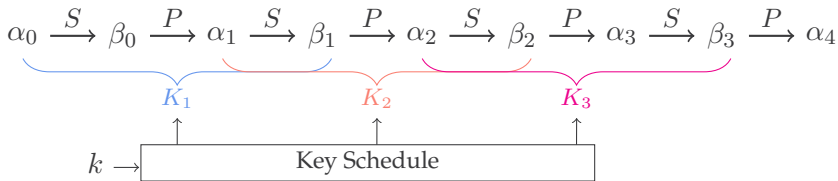
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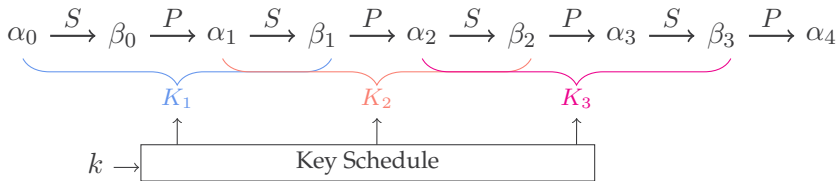
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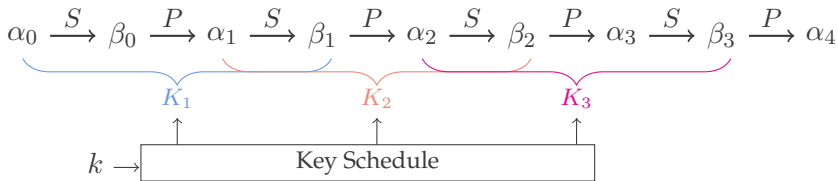
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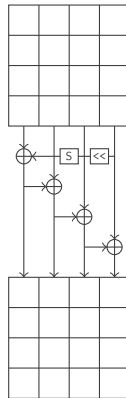


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 - ▶ No solution found \rightarrow singular
 - ▶ Key candidates found \rightarrow Further filter by nonlinear constraints

Singular Characteristics in the AES

Find singular characteristics in AES-128:

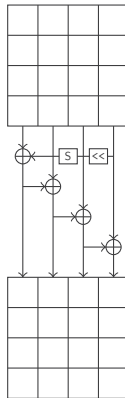


Picture credit:
TikZ for Cryptographers

Singular Characteristics in the AES

Find singular characteristics in AES-128:

- Subspaces of effective keys in every two consecutive rounds

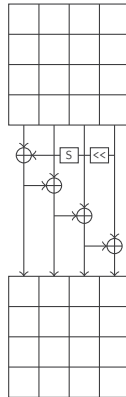


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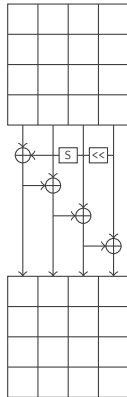


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- 3 out of 4 columns in AES-128 key schedule are linear relations

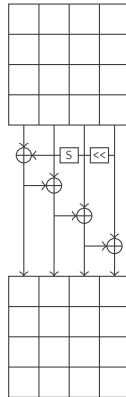


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Singular Characteristics in the AES

Find singular characteristics in AES-128:

- Subspaces of effective keys in every two consecutive rounds
- Build equation systems with key schedule
- 3 out of 4 columns in AES-128 key schedule are linear relations
- Simplify and solve the equation system



Picture credit:
TikZ for Cryptographers

Singular Characteristics in the AES

Examples of 5-round singular characteristics can be found in the AES-128.

$$\begin{array}{ccccccc}
 \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} & \xrightarrow{P} & \begin{pmatrix} 2 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 3 & 0 & 0 & 0 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 3 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 2 & 0 & 0 & 0 \end{pmatrix} \\
 & & & & & & \\
 & \xrightarrow{P} & \begin{pmatrix} 6 & 2 & 1 & 3 \\ 3 & 2 & 3 & 2 \\ 3 & 6 & 2 & 1 \\ 5 & 4 & 1 & 1 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 24 & 27 & 39 & 9d \\ 45 & 36 & 36 & 27 \\ 36 & f1 & 2e & 2d \\ 39 & 2d & 1f & 3a \end{pmatrix} & \xrightarrow{P} & \begin{pmatrix} 6 & 0 & 0 & 0 \\ 0 & 5 & 0 & 0 \\ 0 & 0 & 5 & 0 \\ 0 & 0 & 0 & 36 \end{pmatrix} \\
 & & & & & & \\
 & & \xrightarrow{S} & \begin{pmatrix} e & 0 & 0 & 0 \\ 0 & 9 & 0 & 0 \\ 0 & 0 & d & 0 \\ 0 & 0 & 0 & b \end{pmatrix} & \xrightarrow{P} & \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} .
 \end{array}$$

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MITM attack

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Density of singular characteristics:

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- Enumerate all characteristics given a 3-round differential

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- Enumerate all characteristics given a 3-round differential
- More than 98.47% of all the characteristics are singular
- For the remaining characteristics, we consider the nonlinear constraints from the key schedule and get their effective keys
 - ▶ some of them may also be singular
 - ▶ the number of effective keys is around 2^7 to 2^{10}

Singular Characteristics in the AES

- Different key schedules affect the singularity of a characteristic

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 - ▶ Encrypt a pair of plaintexts under some key with AES-128, track the characteristic

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 - ▶ A valid characteristic in AES-128 is highly probable to be singular in AES-192

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Singular Characteristics in Prince

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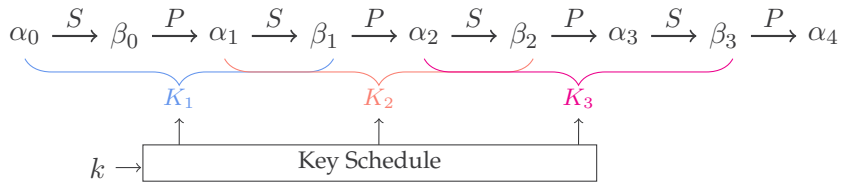
$$\begin{array}{ccccccccc} \begin{pmatrix} 8040 \\ 0000 \\ 4080 \\ 0000 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 8040 \\ 0000 \\ 8040 \\ 0000 \end{pmatrix} & \xrightarrow{M'} & \begin{pmatrix} 8040 \\ 0000 \\ 8040 \\ 0000 \end{pmatrix} & \xrightarrow{SR} & \begin{pmatrix} 8040 \\ 0000 \\ 4080 \\ 0000 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 8050 \\ 0000 \\ 8050 \\ 0000 \end{pmatrix} \\ & & & \xrightarrow{M'} & \begin{pmatrix} 8050 \\ 0000 \\ 8050 \\ 0000 \end{pmatrix} & \xrightarrow{SR} & \begin{pmatrix} 8050 \\ 0000 \\ 5080 \\ 0000 \end{pmatrix} & \xrightarrow{S} & \begin{pmatrix} 2050 \\ 0000 \\ 2050 \\ 0000 \end{pmatrix} \end{array}$$

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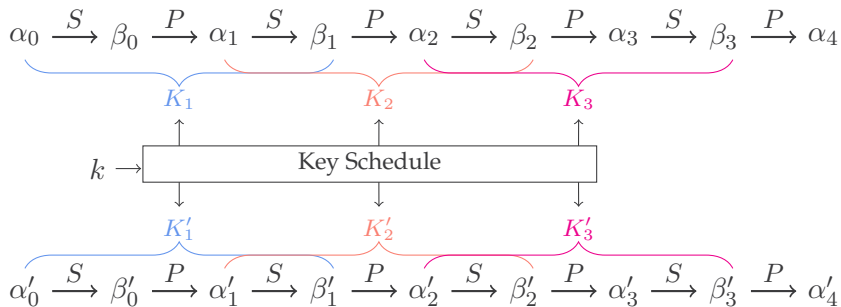
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A 3-round singular characteristic with $\text{EDP} = 2^{-35}$

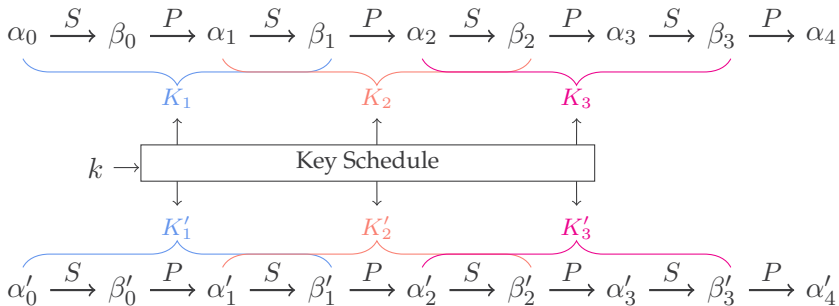
Singular Cluster



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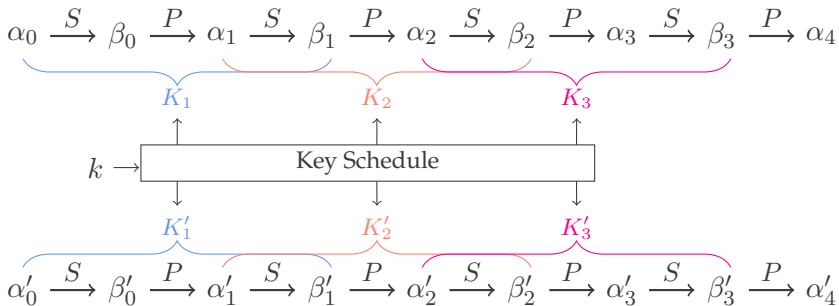


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Differentials/truncated differentials/multiple differentials

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Thank you for your attention!